

PU-003-1164004

Seat No.

M. Sc. (Sem. IV) (CBCS) Examination

August - 2020

CMT - 4004 : Mathematics (Graph Theory)

Faculty Code: 003

Subject Code: 1164004

Time : $2\frac{1}{2}$ Hours] [Total Marks : 70]

Instructions: (1) All questions are compulsory.

(2) Each question carries 14 marks.

1 Answer any seven questions:

 $7 \times 2 = 14$

- (i) Define terms: Degree of a vertex, pendent vertex in graph G and null graph.
- (ii) Define subgraph of a graph G and draw a graph G with its two subgraphs H_1, H_2 so that H_1 and H_2 have five common vertices, but they have no common edges. $(i.e.|V(H_1)\cap V(H_2)|= and \ E(H_1)\cap E(H_2)=\phi).$
- (iii) Define k-regular graph. Also define terms $\delta(G)$ and $\Delta(G)$.
- (iv) Define isomorphism of two graphs. Write down at least two properties for two isomorphic graphs G_1 and G_2 .
- (v) Define Hamiltonian cycle and draw wheel graph W_7 with its Hamiltonian cycle.
- (vi) Define Eulerian graph. Draw a graph G, which admits an Eulerian line and draw another graph H, which can't admit any Eulerian line.
- (vii) Define incidence matrix and write down the incidence matrix for the cycle C_4 .
- (viii) Write down at least three properties of adjacency matrix X(G) for a graph G.

2 Answer any two questions:

 $2 \times 7 = 14$

- (a) Let G be a graph and it contains exactly two odd vertices, say $x, y \in V(G)$. Prove that x and y both lies in the same component of G.
- (b) Let G be a simple graph with n vertices, q edges and k number of components in G. In standard notation prove that $q \le \frac{1}{2}(n-k)(n-k+1)$.
- (c) Let G be a connected graph with $E(G) \neq \phi$. Prove that G is an Eulerian graph if and only if it can be decomposed into edge disjoint cycles.
- (d) State and prove Euler's Theorem.
- 3 Answer any one question:

 $1 \times 14 = 14$

- (a) For a simple connected planner graph G, derive Euler's formula f = e n + 2 and also prove that (i) $e \ge \frac{3f}{2}$ (ii) $e \le 3n 6$. Using these prove that K_5 and $K_{3,3}$ both are non-planner graphs, where e = number of edges in G, n-number of vertices in G and f = number of faces in the planner graph G.
- (b) Let G be a simple graph, |V(G)| > 2 and $d_G(v) \ge \frac{n}{2}$, $\forall v \in V(G)$. Prove that G is a Hamiltonian graph. Also define closure of a graph G and write down closure for K_4, C_5 .
- 4 Answer any two questions:

 $2 \times 7 = 14$

- (a) Define minimally connected graph. Prove that a graph G is a minimally connected graph if and only if it is a tree.
- (b) Let T be a tree with $V(T) \neq \phi$. Prove that T has either one center two centers. In the case it has two centers they must be adjacent in T.
- (c) Prove that a connected graph G, admits a spanning tree.

5 Answer any two questions:

 $2 \times 7 = 14$

- (a) Define weighted graph and minimal spanning tree. Write down two algorithms to obtain minimal spanning tree for a weighted connected graph G, in detail.
- (b) Let G be a connected graph with |V(G)>2. Prove that the vertex connectivity for $G \le$ the edge connectivity for G.
- (c) Define a separable graph. Prove that for a separable graph G, v is a cut vertex in G if and only if there are two vertices $x, y \in V(G) \{v\}$ such that every path in G between x and y passes through v.
- (d) Let T be a tree with $|V(T)| \ge 2$. Prove that T is a 2-chromatic graph.